

A Queue Management Algorithm for Intra-Flow Service Differentiation in the „Best Effort“ Internet

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ICCCN 99, Natick, MA

October 12, 1999



Overview

- Introduction
 - Motivation: graceful degradation under congestion for real-time multimedia flows in the Internet
- Differential RED algorithm (DiffRED)
 - Properties, Differences to RED
- Evaluation
 - Traffic Model
 - Results
- Conclusions

Motivation

- Loss sensitivity of Internet real-time flows:
 - Video: bursty frame losses (packet losses affecting several frames)
 - Voice: bursty packet losses (dependent on codec)
 - „drop-outs“, high perceptual impact
- Solutions:
 - Reservation (IntServ):
complete deployment incl. charging
 - Exploit flow inhomogeneity (Application-layer filtering within the network):
payload-specific, large amount of resources needed, might affect security, difficult to apply to voice

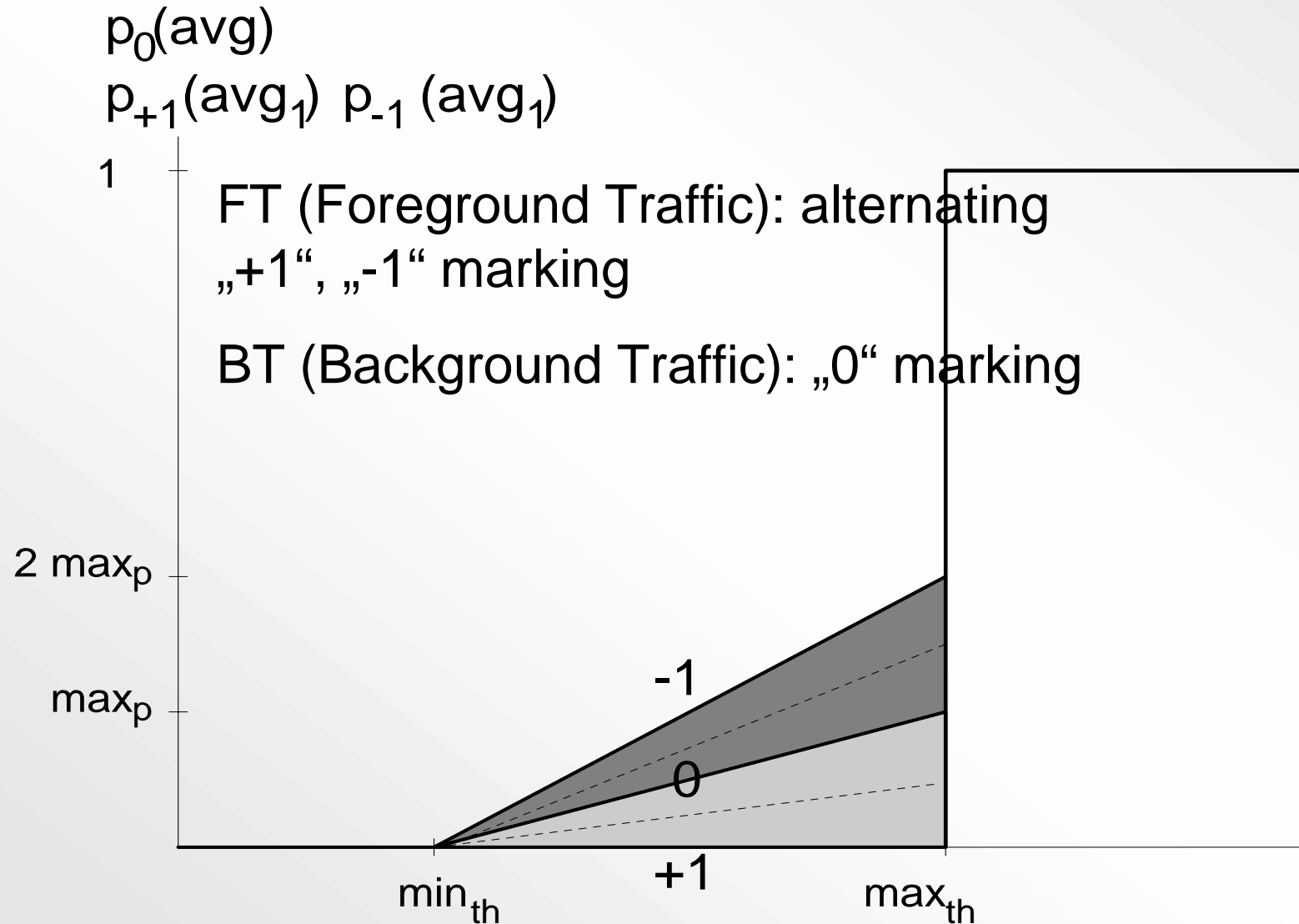
Motivation (cont'd)

- Solutions:
 - End-to-End loss recovery (FEC, loss concealment): efficiency also subject to loss patterns, might worsen congestion
- Adequate mapping of applications' requirements (ADU structure, End-to-End loss recovery) to *simple* network mechanisms
- Provision of a *service* which offers control over (transient) loss distribution / correlation
- Bridge the gap between “Best Effort” and full QoS deployment (QoS migration)

Approach

- What basic mechanisms are needed at a gateway to realize such a service ?
- Simple queue management algorithms (RED) already influence loss correlation (gradual adjustment of the drop probability)
- RIO extends RED to provide *inter-flow* service differentiation
- Application of RIO approach to *intra-flow* service differentiation

Differential RED (DiffRED)



DiffRED: Issues (1)

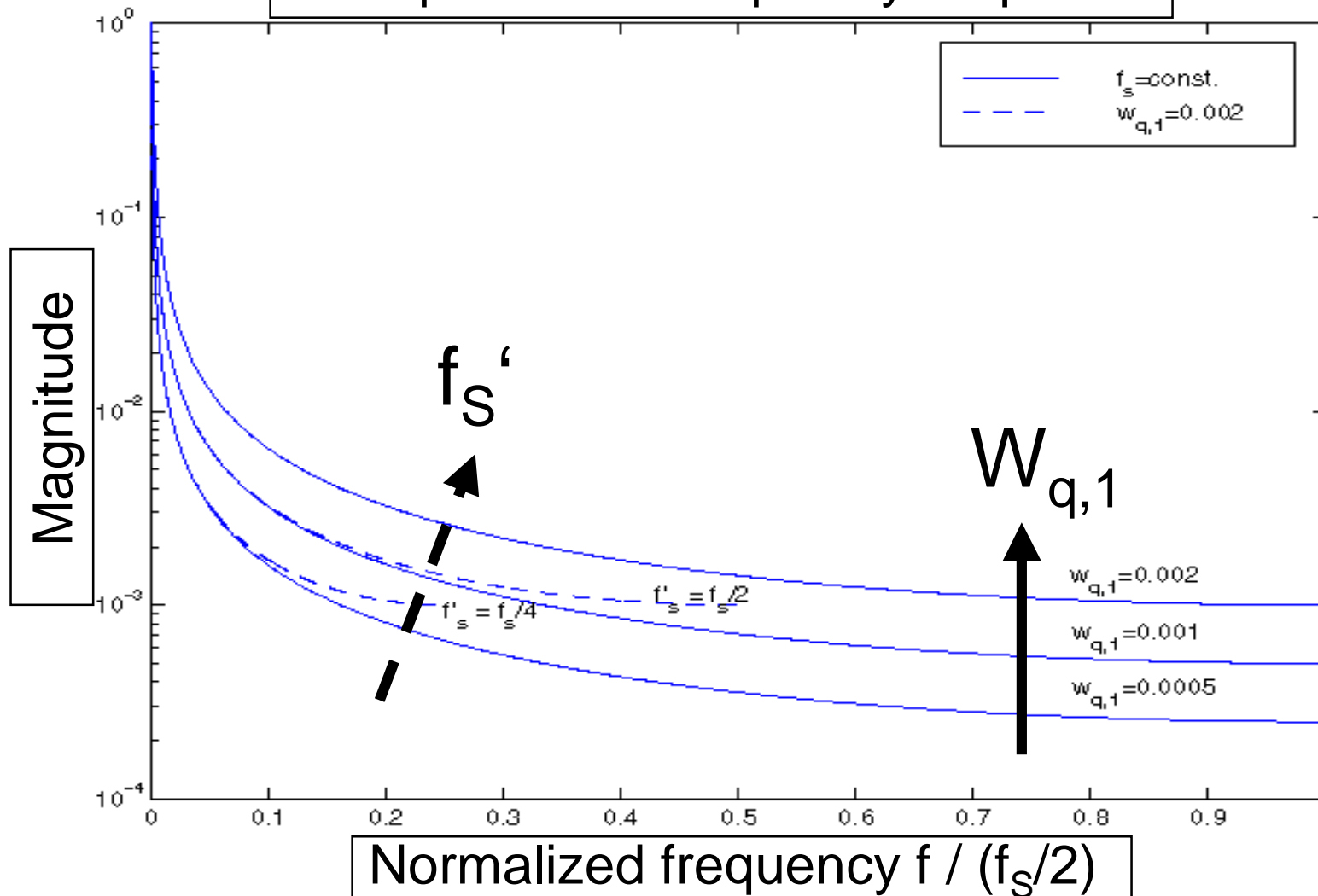
- „Differential“ loss probability curves (compensation of lower/higher FT loss probabilities in the long term)
- Queue state (*avg*) might change substantially between FT arrivals

$$avg \leftarrow (1 - w_q)avg + w_q q$$

- Possible solutions :
 - $w_{q,1} = f(FT, BT \text{ arrivals}) \geq avg_1$
 - $q_1 = f(FT \text{ arrivals})$: sub-sampling of q at FT arrivals
 $\geq avg_1$

DiffRED: Sub-sampling

Low-pass filter frequency response



DiffRED: Issues (2)

- Irregular partition of +1/-1 arrivals
- Possible solutions:
 - monitor and penalize misbehaving flows
 - *adjust loss probability curves to ratio of +1/-1*
- Injection of -1 traffic to mark a flow entirely as +1
- Possible solutions:
 - monitor and penalize misbehaving flows
 - *volume-based charging*

DiffRED: Summary

- „Differential“ loss probability curves
(compensation of lower/higher probabilities in the long term)
- *avg*: sub-sampling of q
- Monitoring of $+1/-1$ arrivals
- fair loss sharing between FT and BT

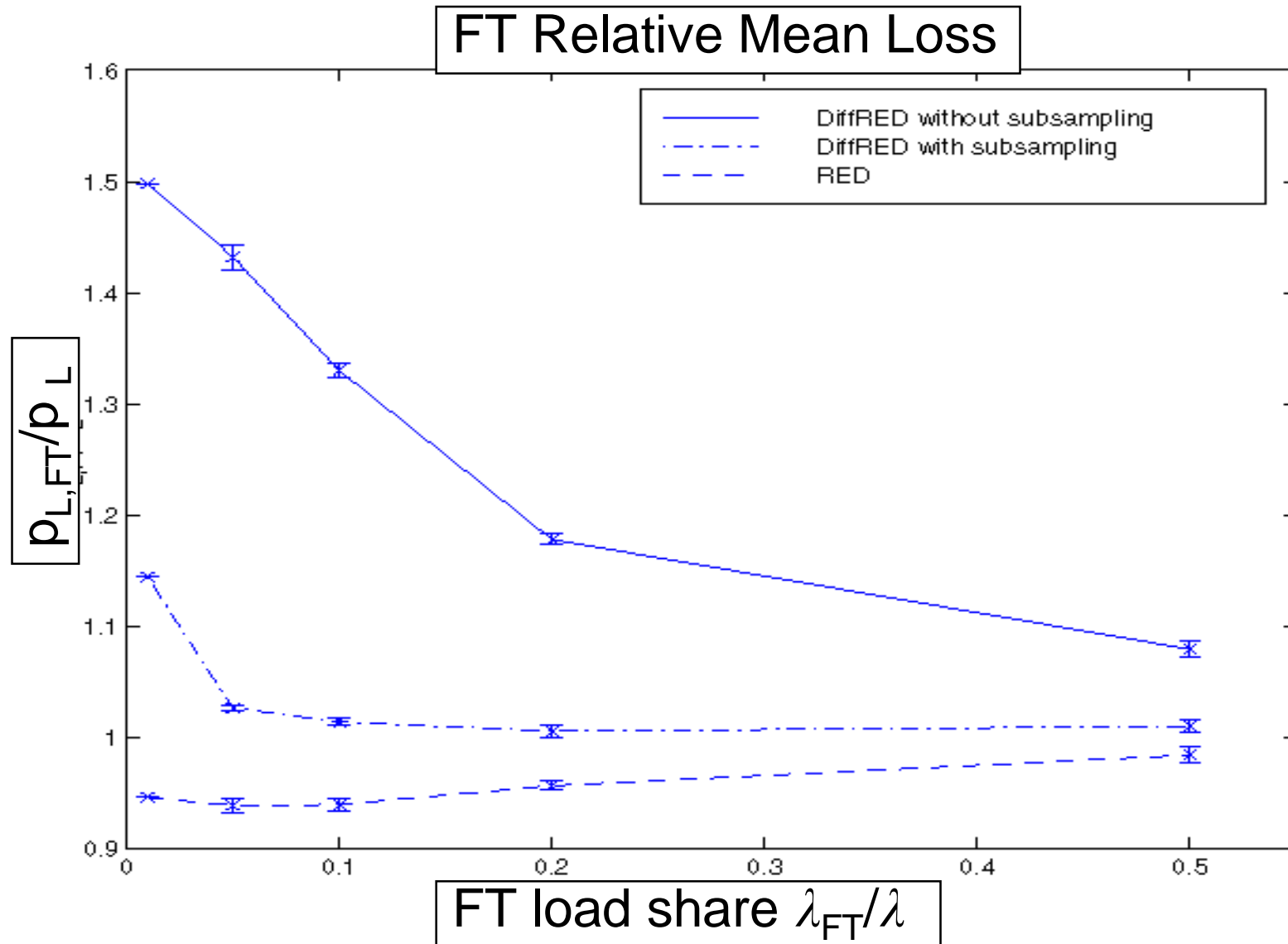
Results

- Traffic Model (based on measurement studies)

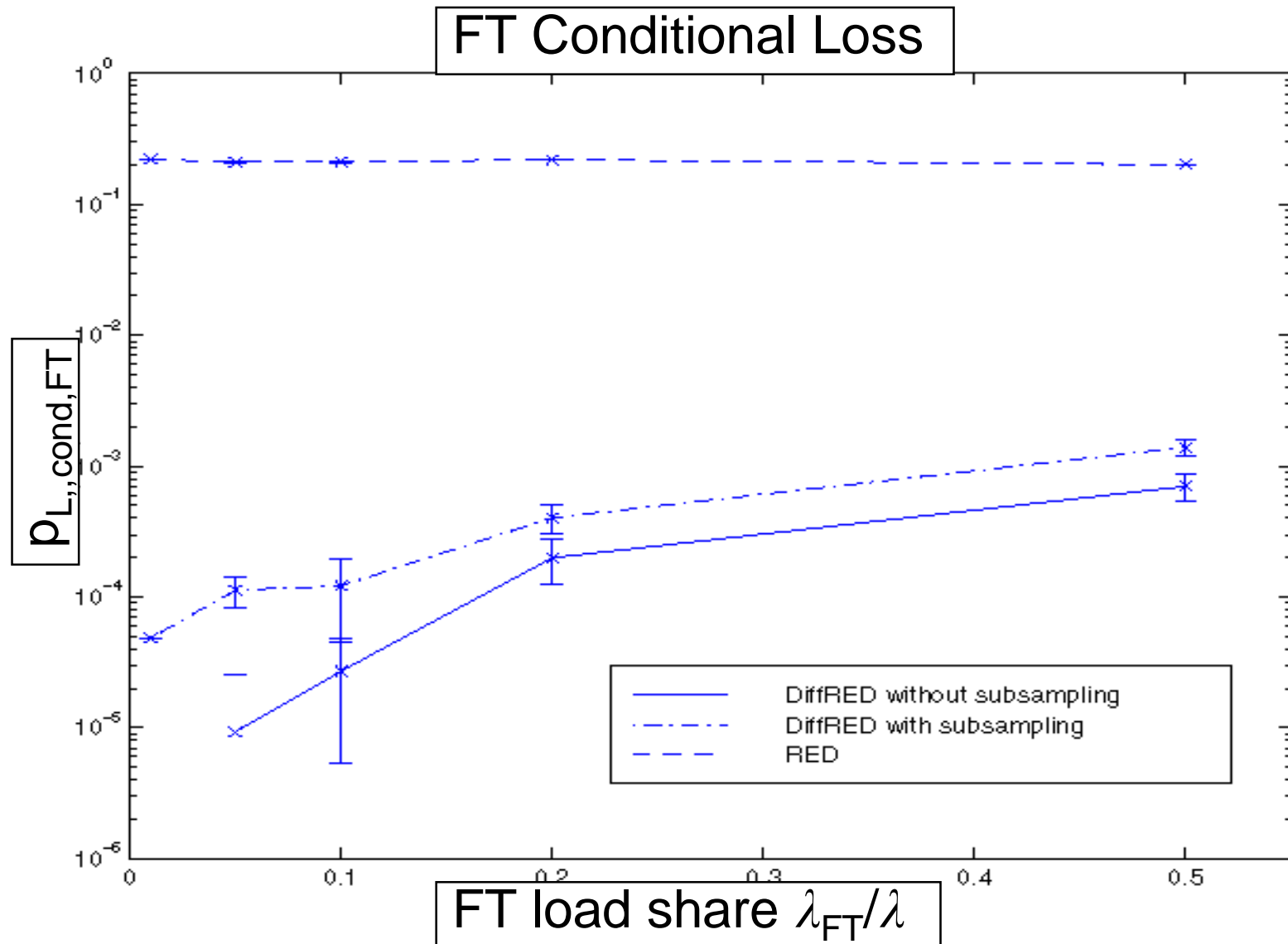
<i>Traffic type</i>	<i>H-BT</i>	<i>D-BT</i>	<i>FT (voice)</i>
flow share (%) of BT	75	25	-
peak bandwidth ($\frac{kBit}{s}$)	256	30...34	83.2
packet size (bytes)	560	128	208
on/off distribution	Pareto	Expo.	Expo.
shape parameter	1.9	—	—
mean burst (packets)	20	4	18
mean ontime (s)	0.35	0.12...0.14	0.36
mean offtime (s)	0.7	0.12...0.14	0.64

- Experiment: variation of the FT load share λ_{FT}/λ at a fixed traffic intensity $\rho=\lambda/\mu=0.95$, single gateway

Results



Results



Conclusions

- Control loss distribution for certain flows while maintaining RED properties in the long term
 - applications' requirements (ADU, e2e loss recovery) can be mapped on a simple network mechanism
- No complete QoS architecture needed, partial deployment beneficial, suitable framework: DiffServ AF (three drop precedences within a class)
- Further study for other (bursty) traffic types needed